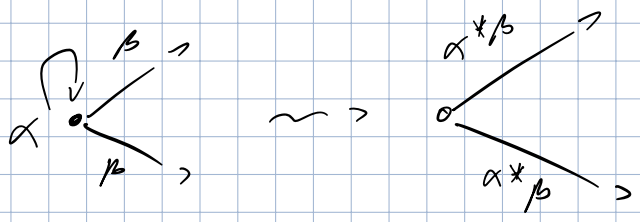
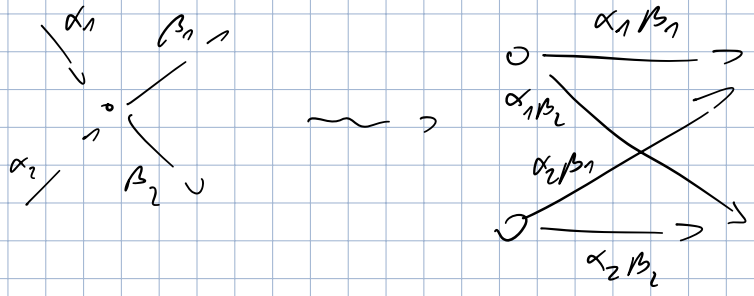
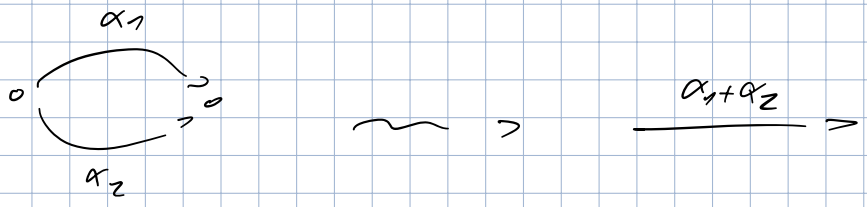
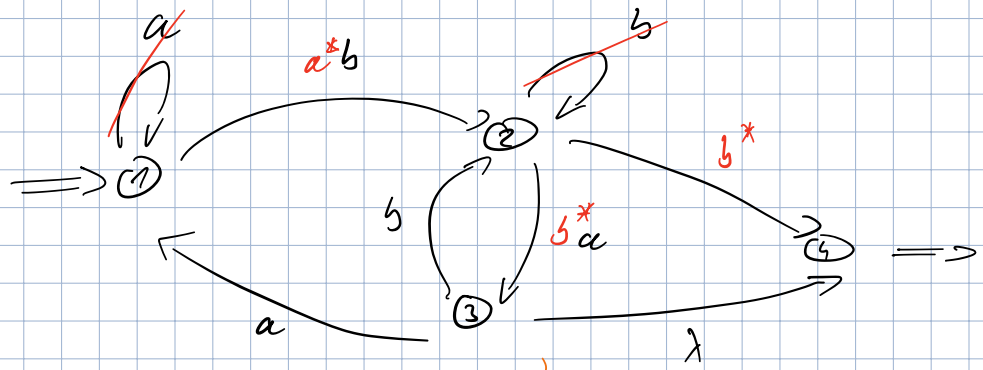
Generování automatu:

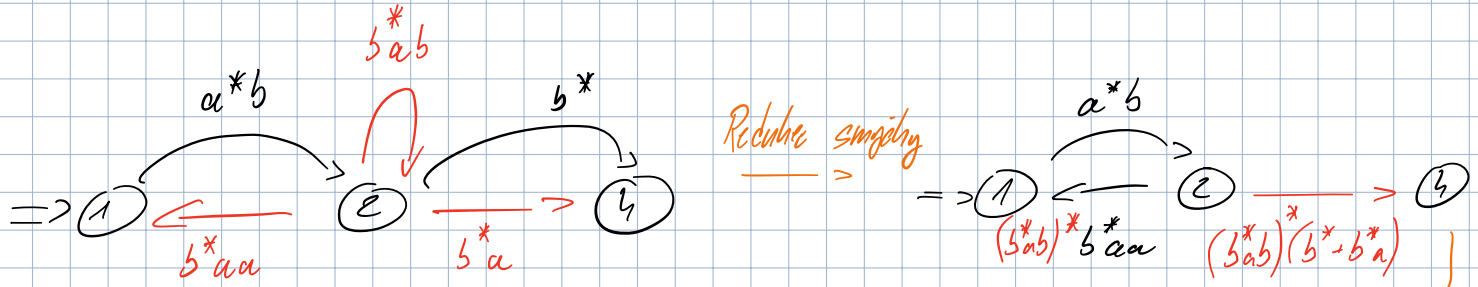




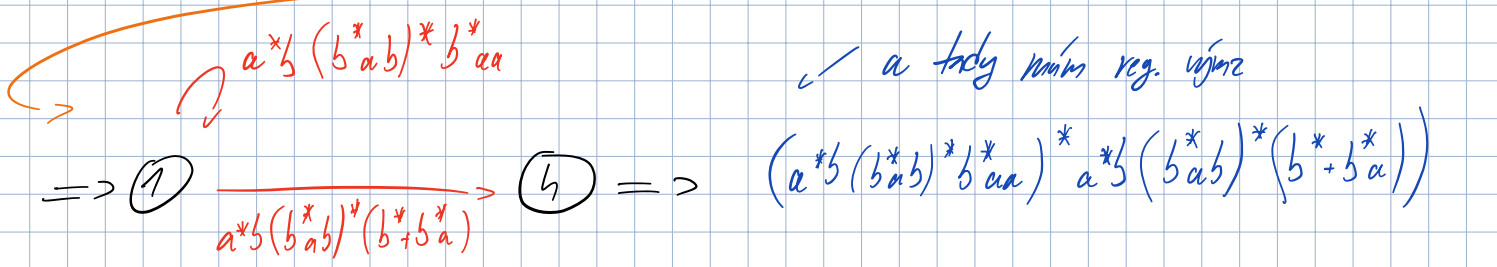




Redukce 3



Zjednodušení ① — ②



a tady máme reg. výraz

$$\left(a^*b\left(b^*ab\right)^*b^*aa\right)^*a^*b\left(b^*ab\right)^*\left(b^*+b^*a\right)$$

$\{ w \in \{a,b\}^* \mid w$ začínají na $\underline{ba}$ a končí na $\underline{ab} \}$

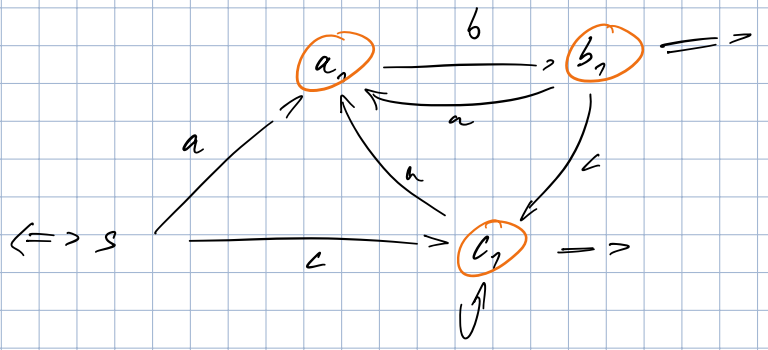
$ba(a+b)^*ab + bab$

$\{ w \in \{0,1\}^* \mid |w|_0 = 2i,\ i \in \mathbb{N} \}$ —> tedy sudý počet nul

$(1^*01^*01^*)^*$

Dostaneme Reg —> chceme automat

$(a_1b_1 + c_1)^*$

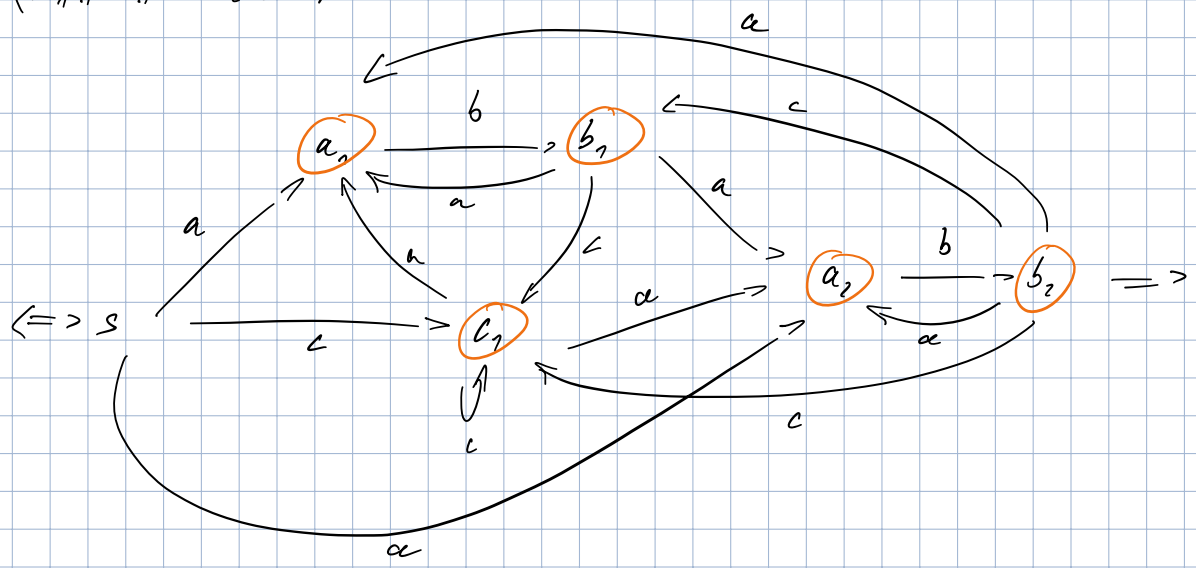


1) začínají
2) navazují po sobě
3) poslechni
4) λ ?

c

↳ tohle je jednodušší než stavba přes triviální automaty

$((a_1b_1 + c_1)^*(a_2b_2)^+)^*$

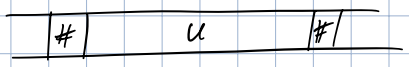


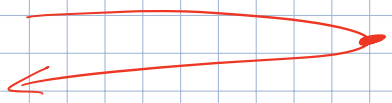
2DA → dvoucestný lineární automat: → mám 1DA $A = (Q, \Sigma, \delta, q_0, F)$

reg.

$\{ \# u \# \mid uu^R \in L(A) \}$

→ mám omezenou pamět, ale neomezený čas





→ stále přijímá jen reg. jazyky

DÚ: Sestrojte tento automat, pomocí metody který sestrojíte 1-cestný.