# ? C++ - parallelization and synchronization

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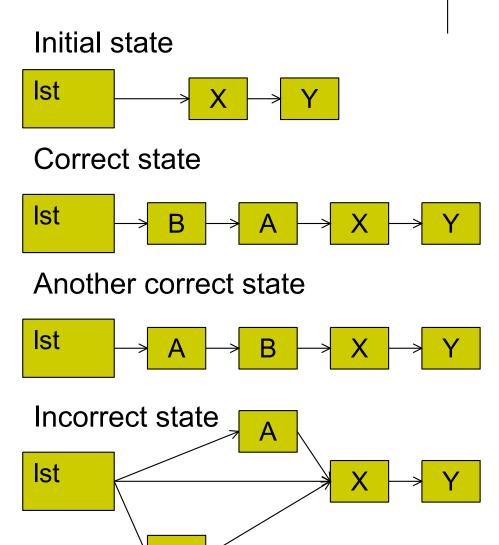
#### ? The problem

- ? Race conditions
  - Separate threads with shared state
  - Result of computation depends on OS scheduling

# ? Race conditions – simple demo



- ? Linked list
- ? Shared state
  List Ist;
- ? Thread A
  Ist.push\_front(A);
- ? Thread B lst.push\_front(B);



# ? Race conditions – advanced demo



```
struct Counter {
                           ? Shared state
  Counter():value(0) { }
                           Counter c;
  int value;
                           ? Thread A
  void increment()
                           c.increment();
  { ++value; }
                           cout << c.get();
  void decrement()
    --value; }
                           ? Thread B
  int get()
                           c.increment();
  { return value; }
                           cout << c.get();
                           Possible outputs
                           12, 21, 11
```

#### ? C++ features

- ? C++ 11
  - Atomic operations
  - ! Low-level threads
  - High-level futures
  - Synchronization primitives
  - ? Thread-local storage

- ? C++14 features
  - Shared timed mutex
- ? C++17 features
  - Parallel algorithms
  - ? Shared mutex
- ? C++20 features
  - Stop tokens
  - Semaphore
  - ? Coordination types



- ? Low-level threads
  - ? Header <thread>
  - ? thread class
  - Pork-join paradigm
  - ? Namespace this\_thread



- ? Class thread
  - ? Constructor
    - ! template <class F, class ...Args>
      explicit thread(F&& f, Args&&... args);
  - ? Destructor
    - If joinable() then terminate()
  - bool joinable() const noexcept;
  - ? void join();
    - Blocks, until the thread \*this has completed
  - ? void detach();
  - ② id get\_id() const noexcept;
  - ② static unsigned hardware\_concurrency();



- Class jthread
  - Like thread, autostops+autojoins on destruction
  - Provides stop token
    - Internal member of std::stop\_source type
    - Constructor accepts function with std::stop\_token as first argument
    - Proposition of the property of the property
    - Interface functions get\_stop\_source, get\_stop\_token, and request\_stop



- ? Namespace this\_thread
  - 1 thread::id get\_id() noexcept;
    - Unique ID of the current thread
  - void yield() noexcept;
    - Opportunity to reschedule
  - sleep\_for, sleep\_until
    - Blocks the thread for relative/absolute timeout



```
Demo
#include <iostream>
#include <thread>
void thread_fn() { std::cout << "Hello from thread" <<</pre>
  std::endl; }
int main(int argc, char **argv) {
  std::thread thr(&thread_fn);
  std::cout << "Hello from main" << std::endl;
  thr.join();
  return 0;
```



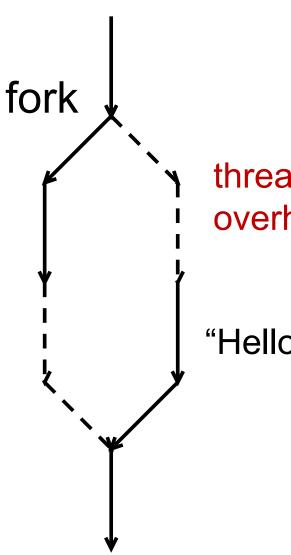
in" "Hello from thread" join

"Hello from main"



"Hello from main"

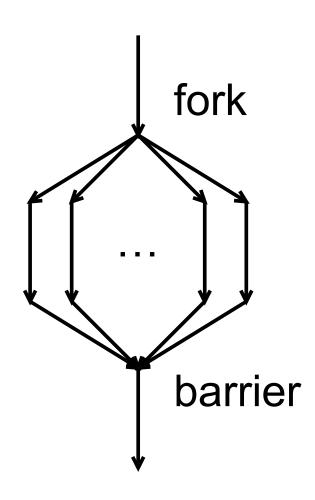
blocked on join



thread creation overhead

"Hello from thread"







#### Demo #include <iostream> #include <thread> #include <vector> int main(int argc, char \*\*argv) { std::vector<std::thread> workers; for(int i=0;i<10;++i) workers.push\_back(std::thread([i]() { std::cout << "Hello from thread " << i << std::endl; std::cout << "Hello from main" << std::endl; for(auto &t : workers) t.join(); return 0;



- Passing arguments to threads
  - By value
    - Safe, but you MUST make deep copy
  - By move (rvalue reference)
    - Safe, as long as strict (deep) adherence to move semantics
  - By const reference
    - Safe, as long as object is guaranteed deep-immutable
  - By non-const reference
    - Safe, as long as the object is monitor



- ? Futures
  - ? Header <future>
  - ! High-level asynchronous execution
  - ? Future
  - ? Promise
  - ? Async
  - Error handling



- ? Shared state
  - Consist of
    - Some state information and some (possibly not yet evaluated) result, which can be a (possibly void) value or an exception
  - Asynchronous return object
    - Object, that reads results from an shared state
  - Waiting function
    - Potentially blocks to wait for the shared state to be made ready
  - Asynchronous provider
    - Object that provides a result to a shared state



- ? Future
  - ? std::future<T>
  - Puture value of type T
  - Retrieve value via get()
    - Waits until the shared state is ready
  - wait(), wait\_for(), wait\_until()

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- ? valid()
- ? std::shared\_future<T>
  - Value can be read by more then one thread



- ? Async
  - ? std::async
  - ! Higher-level convenience utility
  - Launches a function potentially in a new thread
- ? Async usage

```
int foo(double, char, bool);
auto fut = std::async(foo, 1.5, 'x', false);
auto res = fut.get(); > toble ale unotation; i senoue volain...

> toble hade std:: fature (int) > tady to yunda, valid() ut potom
hade == false.
```



- Packaged task
  - std::packaged\_task
  - How to implement async with more control
  - Wraps a function and provides a future for the function result value, but the object itself is callable



Packaged task usage
std::packaged\_task<int(double, char, bool)>
 tsk(foo);
auto fut = tsk.get\_future();
std::thread thr(std::move(tsk), 1.5, 'x', false);
auto res = fut.get();



- ? Promise
  - ? std::promise<T>
  - ? Lowest-level
  - Steps
    - Calling thread makes a promise
    - Calling thread obtains a future from the promise
    - The promise, along with function arguments, are moved into a separate thread
    - ? The new thread executes the function and fulfills the promise
    - The original thread retrieves the result



- ? Promise usage
  - ? Thread A

```
std::promise<int> prm;
auto fut = prm.get_future();
std::thread thr(thr_fnc, std::move(prm));
auto res = fut.get();
? Thread B
void thr_fnc(std::promise<int> &&prm) {
    prm.set_value(123);
}
```



#### ? Constraints

- A default-constructed promise is inactive
  - ? Can die without consequence
- A promise becomes active, when a future is obtained via get\_future()
  - Only one future may be obtained
- ? A promise must either be satisfied via Set\_value(), or have an exception set via Set\_exception()
  - A satisfied promise can die without consequence
  - get() becomes available on the future
  - ? A promise with an exception will raise the stored exception upon call of get() on the future
  - ? A promise with neither value nor exception will raise "broken promise" exception



- ? Exceptions
  - All exceptions of type std::future\_error
    - ? Has error code with enum type Std::future\_errc

```
inactive promise
std::promise<int> pr;
fine, no problem

inactive promise

inactive promise
inactive promise

inactive promise

inactive promise

std::promise<int> pr;

auto fut1 = pr.get_future();

auto fut2 = pr.get_future();

inactive promise

std::promise<int> pr;

auto fut1 = pr.get_future();

inactive promise
```

? active promise, unused error "Future already retrieved" std::promise<int> pr;

auto fut = pr.get\_future();
// fine, no problem
// fut.get() blocks indefinitely



```
? satisfied promise
std::promise<int> pr;
auto fut = pr.get_future();
{ std::promise<int>
pr2(std::move(pr));
  pr2.set_value(10);
auto r = fut.get();
// fine, return 10
```

```
? too much satisfaction
std::promise<int> pr;
auto fut = pr.get_future();
{ std::promise<int>
pr2(std::move(pr));
  pr2.set_value(10);
  pr2.set_value(11);
// error "Promise already
satisfied"
auto r = fut.get();
```



```
? exception
std::promise<int> pr;
auto fut = pr.get_future();
{ std::promise<int> pr2(std::move(pr));
  pr2.set_exception(
    std::make_exception_ptr(
      std::runtime_error("bububu")));
auto r = fut.get();
// throws the runtime error
```



```
Proken promise
std::promise<int> pr;
auto fut = pr.get_future();
{ std::promise<int> pr2(std::move(pr));
    // error "Broken promise"
}
auto r = fut.get();
```



### ? Synchronization primitives

- ? Synchronization primitives
  - Mutual exclusion
    - Peaders <mutex> and <shared\_mutex>
  - Condition variables
    - ? Header <condition\_variable>
  - Semaphore
    - ? Header <semaphore>



#### ? Mutex

#### ? Mutex

- A synchronization primitive that can be used to protect shared data from being simultaneously accessed by multiple threads
- mutex offers exclusive, non-recursive ownership semantics
  - ? A calling thread owns a mutex from the time that it successfully calls either lock or try\_lock until it calls unlock
  - ? When a thread owns a mutex, all other threads will block (for calls to lock) or receive a false return value (for try\_lock) if they attempt to claim ownership of the mutex
  - A calling thread must not own the mutex prior to calling lock or try\_lock
- The behavior of a program is undefined if a mutex is destroyed while still owned by some thread



#### ? Mutex example

```
? Shared state
List Ist;
std::mutex mtx;
? Thread A
mtx.lock();
lst.push_front(A);
mtx.unlock();
? Thread B
mtx.lock();
lst.push_front(B);
mtx.unlock();
```



#### ? Mutex variants

- Other mutex variants
  - ? timed\_mutex
    - In addition, timed\_mutex provides the ability to attempt to claim ownership of a timed\_mutex with a timeout via the try\_lock\_for and try\_lock\_until
  - ? recursive\_mutex
    - ? exclusive, recursive ownership semantics
      - A calling thread owns a recursive\_mutex for a period of time that starts when it successfully calls either lock or try\_lock. During this period, the thread may make additional calls to lock or try\_lock. The period of ownership ends when the thread makes a matching number of calls to unlock
      - When a thread owns a **recursive\_mutex**, all other threads will block (for calls to **lock**) or receive a false return value (for **try\_lock**) if they attempt to claim ownership of the **recursive\_mutex**
      - The maximum number of times that a **recursive\_mutex** may be locked is unspecified, but after that number is reached, calls to **lock** will throw **std::system\_error** and calls to **try\_lock** will return false
  - ? recursive\_timed\_mutex
    - ? Combination



#### ? Mutex variants

- Other mutex variants
  - ? std::shared\_mutex
    - ? Additionally multiple threads can make shared lock using lock\_shared()
    - Either exclusive lock or shared lock
  - ? std::shared\_timed\_mutex



#### ? Mutex wrappers

- 1 std::unique\_lock
  - ! Lock class with more features
    - Timed wait, deferred lock
- ② std::lock\_guard
  - Scope based lock (RAII)
  - Linked list demo, code for one thread

```
std::lock_guard<std::mutex> lk(mtx);
lst.push_front(X);
}
```



#### ? Mutex wrappers and others

- ? Shared lock wrapper
  - std::shared\_lock
    - Calls lock\_shared for the given shared mutex
- ? Variadic wrapper
  - ? template <typename ... MutexTypes> class
    scoped\_lock;
    - Multiple locks at once, RAII, deadlock avoidance
- ? Interference size
  - std::size\_t
     hardware\_destructive\_interference\_size;
  - Size of a cache line



#### ? Locking algorithms

- ? std::lock
  - locks specified mutexes, blocks if any are unavailable, deadlock avoidance
- ? std::try\_lock
  - attempts to obtain ownership of mutexes via repeated calls to try\_lock

```
// don't actually take the locks yet
std::unique_lock<std::mutex> lock1(mtx1, std::defer_lock);
std::unique_lock<std::mutex> lock2(mtx2, std::defer_lock);
// lock both unique_locks without deadlock
std::lock(lock1, lock2);
```



#### ? Call once

- 1 std::once\_flag
  - ? Helper object for std::call\_once
- ? std::call\_once
  - invokes a function only once even if called from multiple threads

```
std::once_flag flag;
void do_once() {
    std::call_once(flag, [](){ do something only once }); }
std::thread t1(do_once);
std::thread t2(do_once);
```



#### ? Condition variable

#### ② std::condition\_variable

- Can be used to block a thread, or multiple threads at the same time, until
  - ? a notification is received from another thread
  - a timeout expires, or
  - ? a spurious wakeup occurs
    - Appears to be signaled, although the condition is not valid
    - Verify the condition after the thread has finished waiting
- ? Works with std::unique\_lock
- wait atomically manipulates mutex, notify does nothing



#### ? Condition variable example

```
std::mutex m;
 std::condition_variable cond_var;
 bool done = false; bool notified = false;
? Producer
                                       ? Consumer
for () {
                                       std::unique_lock<std::mutex> lock(m);
                                      while(!done) {
 // produce something
                                         while (!notified) {
 { std::lock_guard<std::mutex>
   lock(m);
                                          // loop to avoid spurious wakeups
    queue.push(item);
                                           cond var.wait(lock);
    notified = true; }
  cond_var.notify_one();
                                         while(!queue.empty()) {
                                           queue.pop();
std::lock quard<std::mutex> lock(m);
                                          // consume
notified = true;
done = true;
                                         notified = false;
cond_var.notify_one();
```



#### ? Semaphore

- ? Counting semaphore
  - std::counting\_semaphore
  - Constructor sets the count
  - Manipulation
    - ? acquire(), release(count=1)
- ? Binary semaphore
  - std::binary\_semaphore



## ? Coordination types

- ? Latches
  - ? Header < latch >
  - ? Thread coordination mechanism
    - Block threads until an expected number of threads arrive
  - Single use
- ? Barriers
  - ? Header <barrier>
  - Sequence of phases
    - Each call to arrive() decrements expected count, the thread can then wait()
    - When count==0, the completion function is called and all blocked threads are unblocked
    - Expected count is reset to the previous value
  - ? Constructor
    - barrier(ptrdiff\_t expected, CompletionFunction f)



#### ? Stop tokens

- ? Asynchronously request to stop execution of an operation
  - Shared state among associated stop\_source, stop\_token, and stop\_callback
  - ? stop\_token
    - Interface for querying whether a stop request has been made or can ever been made
      - bool stop\_requested(), bool stop\_possible()
  - ? stop\_source
    - !mplements the semantics of making stop request
    - ? Creates stop\_tokens
      - ? stop\_token get\_token()
    - Makes stop request
      - ? request\_stop()
  - ? stop\_callback
    - Invokes callback function when stop request has been made



## ? Thread-local storage

- ? Thread-local storage
  - Added a new storage-class
  - Use keyword thread\_local
    - Must be present in all declarations of a variable
    - Only for namespace or block scope variables and to the names of static data members
      - ? For block scope variables Static is implied
  - Storage of a variable lasts for the duration of a thread in which it is created



## ? Parallel algorithms

- ? Parallelism
  - In headers <algorithm>, <numeric>
  - Parallel algorithms
    - Execution policy in <execution>
      - seq execute sequentially
      - par execute in parallel on multiple threads
      - par\_unseq execute in parallel on multiple threads, interleave individual iterations within a single thread, no locks
      - unseq execute in single thread+vectorized
    - ? for\_each
    - reduce, scan, transform\_reduce, transform\_scan
      - Inclusive scan like partial\_sum, includes i-th input element in the i-th sum
      - Exclusive scan like partial\_sum, excludes i-th input element from the i-th sum
    - No exceptions should be thrown
      - ? Terminate



## ? Parallel algorithms

- Parallel algorithms
  - Not all algorithms have parallel version
  - adjacent\_difference, adjacent\_find, all\_of, any\_of, copy\_if, copy\_n, count, if, equal, exclusive\_scan, fill\_n, find, find\_end, find\_first\_of, find\_if, find\_if\_not, for each, for each\_n, generate, generate\_n, includes, inclusive\_scan, inner\_product, inplace\_merge, is\_heap, is\_heap\_until, is\_partitioned, is\_sorted, is\_sorted\_until, lexicographical\_compare, max\_element, merge, min\_element, minmax\_element, mismatch, move, none\_of, nth\_element, partial\_sort, partial\_sort\_copy, partition, partition\_copy, reduce, remove, remove\_copy, remove\_copy\_if, remove\_if, replace, replace\_copy, replace\_copy\_if, replace\_if, reverse, reverse\_copy, rotate\_copy, search, search\_n, set\_difference, set\_intersection, set\_symmetric\_difference, set\_union, sort, stable\_partition, stable\_sort, swap\_ranges, transform\_transform\_evaluative\_copy, partition, transform, transform\_exclusive\_scan, transform\_inclusive\_scan, transform\_reduce, uninitialized\_copy, uninitialized\_copy\_n, uninitialized\_fill, uninitialized\_fill\_n, unique, unique\_copy



#### ? C++ extension – executors

- ? Executors
  - Now separate TS, not finished in C++23 timeframe, maybe in C++26?
- ? Executor
  - Controls how a task (=function) is executed
  - ? Direction
    - ? One-way execution
      - ? Does not return a result
    - ? Two-way execution
      - ? Returns future
    - ? Then
      - Execution agent begins execution after a given future becomes ready, returns future
  - ? Cardinality
    - Single
      - ? One execution agent
    - Bulk executions
      - ? Group of execution agents
      - ? Agents return a factory
- ? Thread pool
  - Controls where the task is executed

# ? C++ extensions – concurrency



- ? Concurrency
  - TS published, depends on executors TS
  - Improvements to future
    - future<T2> then(F &&f)
      - Execute asynchronously a function f when the future is ready

# ? C++ extension – transactional memory



- TS v1 finished, never used
- ? TS v2 in progress
- ? Transactional memory
  - ? Atomic blocks
    - Transactional behavior
    - Exception inside the block leads to undefined behavior

```
unsigned int f()
{
    static unsigned int i = 0;
    atomic do {
        ++i;
        return i;
    }
}
```