

Bayesian model consists of:

1) Single par. mechanism (x, p)

- n bidders, with valuations $v_1 - v_n$

- $x = \text{allocation rule}$

- $p = \text{payment rule}$

- feasible outcomes X

2) Valuations $v_1 - v_n$ are drawn from independent prob. distr. $F_1 - F_n$

with densities $f_1 - f_n$ with support on $\langle 0, v_{\max} \rangle$

$$F_i(z) = P_v [v_i \leq z]$$

$$\cdot f_i(x) = \frac{d}{dx} F_i(x), \quad F_i(x) = \int_0^x f_i(z) dz$$

3) $F_1 - F_n$ are known to the seller

- we try to maximize $E[\sum_i p_i(v)]$ = expected revenue

$$v \sim (F_1 - F_n)$$

$$- E[X(v)] = \int_0^{v_{\max}} X(v) f(v) dv$$

Example 1: (single-item, 1 bidder)

- consider uniform prob. dist. F on $\langle 0, 1 \rangle$ $F(x) = x \forall x \in \langle 0, 1 \rangle$

\Rightarrow exp. revenue is $r \cdot (1 - F(r)) = r \cdot (1 - r) \Rightarrow$ this gives maximum $\frac{1}{4}$, for $r = \frac{1}{2}$

Example 2: (single-item, 2 bidders)

- we can apply Vickrey's auction \Rightarrow exp. revenue is $\frac{1}{3}$

- \exists better auction that gives $\frac{5}{12}$.

Thm:

In a single-par env. DSIC (x, p) with prob. distr. $F_1 - F_n$ with densities $f_1 - f_n$

if $F = F_1 \times \dots \times F_n$, then $E[\sum_i p_i(v)] = E[\sum_i \psi_i(v_i) \cdot x_i(v)]$

where $\psi_i(v_i) = v_i - \frac{1 - F_i(v_i)}{f_i(v_i)}$

this would be social surplus without ψ

inevitable loss for not knowing v_i exactly. Therefore it is called virtual social surplus.

- can be negative (if F uniform on $\langle 0, 1 \rangle \Rightarrow \psi_i(v_i) = v_i - \frac{1 - v_i}{1} = 2v_i - 1$)

$\ll -1, 1$ then

Ob: (x, p) is DSIC $\Rightarrow b_i = v_i$ "bids are valuations"

$$\Rightarrow \text{Myerson's lemma} \Rightarrow p_i(v) = \int_0^{v_{\max}} 2 \cdot \frac{\partial}{\partial z} x_i(z, v_{-i}) dz$$

fix i, v_{-i} :

$$\mathbb{E}_{V_i \sim F_i} [p_i(v_i, v_{-i})] = \int_0^{v_{\max}} p_i(v_i, v_{-i}) \cdot f_i(v_i) dv_i = \int_0^{v_{\max}} \int_0^{v_{\max}} 2 \cdot \frac{\partial}{\partial z} x_i(z, v_{-i}) \cdot f_i(v_i) dz dv_i =$$

definition of \mathbb{E}

$$= \int_0^{v_{\max}} \left(\int_0^{v_{\max}} f_i(v_i) dv_i \right) \cdot 2 \cdot \frac{\partial}{\partial z} x_i(z, v_{-i}) dz = \int_0^{v_{\max}} (1 - F_i(z)) \cdot 2 \cdot \frac{\partial}{\partial z} x_i(z, v_{-i}) dz$$

$1 - F_i(z)$

Fabini theorem

$$\text{per partes } \int f g = f \cdot g - \int f' \cdot g : f(z) = (1 - F_i(z))$$

$$g(z) \frac{\partial}{\partial z} x_i(z, v_{-i})$$

$$= \left[(1 - F_i(z)) \cdot 2 \cdot x_i(z, v_{-i}) \right]_0^{v_{\max}} - \int_0^{v_{\max}} (1 - F_i(z) - 2 \cdot f_i(z)) \cdot x_i(z, v_{-i}) dz$$

$= 0$

definition of \mathbb{E}

$$= \int_0^{v_{\max}} \left(2 - \frac{1 - F_i(z)}{f_i(z)} \right) \cdot f_i(z) \cdot x_i(z, v_{-i}) dz = \mathbb{E}_{V_i \sim F_i} [\ell_i(v_i) \cdot x_i(v_i, v_{-i})]$$

only rewriting

$$\Rightarrow \forall i \nexists v_{-i} : \mathbb{E}_{V_i \sim F_i} [p_i(v_i, v_{-i})] = \mathbb{E}_{V_i \sim F_i} [\ell_i(v_i) x_i(v_i, v_{-i})]$$

integration over all v_{-i}

$$\Rightarrow \mathbb{E}_{V \sim F} [p_i(v)] = \mathbb{E}_{V \sim F} [\ell_i(v_i) x_i(v)]$$

$$\Rightarrow \mathbb{E}_{V \sim F} [\mathbb{E}_{V \sim F} [p_i(v)]] = \sum_i \mathbb{E}_{V \sim F} [p_i(v)] = \sum_i \mathbb{E}_{V \sim F} [\ell_i(v_i) x_i(v)] = \mathbb{E}_{V \sim F} [\sum_i \ell_i(v_i) x_i(v)]$$

Let's apply the thm. Consider single-item auction, $F=F_1=\dots=F_n \Rightarrow f=f_1=\dots=f_n$

Assume that F is regular (ℓ is strictly increasing)
(true for F uniform on $(0, 1)$)
 $\Rightarrow \ell=\ell_1=\dots=\ell_n$

Vickrey's auction with reverse price $m \rightarrow$ This is eBay auction!

- winner is the highest bidder, but if bids are $< m$, then there is no winner. Otherwise we get negative $\ell_i(v_i)$
- winner (if any) pays 2nd highest bid or m , whatever is larger.

Under our assumptions, Vickrey's auction with reserve price $\ell^{-1}(0)$ maximizes revenue. \rightarrow since F is regular, only one image exists.

Dh: We want to minimize expected virtual valuation

$$X = \{x_1 - x_n \in \{0, 1\}^n : \sum x_i \leq 1\} \Rightarrow \text{winner should have highest } \ell_i(v_i)$$

- if $\ell_1(v_1), \dots, \ell_n(v_n) < 0 \Rightarrow \text{no winner}$

- Myerson's Lemma gives the payment rule - pay $m = \ell^{-1}(0)$ or
2nd highest $\ell_i(v_i)$

\Rightarrow Vickrey's auction with reserve price $\ell^{-1}(0)$. ☒

\hookrightarrow allocation rule is monotone (if F regular)