

No regret dynamics

$G = (P, A, C)$ of n players

$C \rightarrow$ cost function: -1. utility

$C \in (-1, 1)$ - if not, we can scale...

Play G T -times, the player $i \sim i$ -th agent chooses p^t as his mixed strategy at step t .

\rightarrow prob. distr. on A_i obtained by RW alg. with average regret $\leq \epsilon$

Losses of i are captured by C_i

$$\text{for } a_i \in A_i: \ell^t(a_i) = \mathbb{E}_{a_{-i} \sim p_{-i}^t} [C_i(a_i, a_{-i}^t)]$$

\rightarrow all other play according to the RW alg. strategy

Application:

The expected cost $C_2(s) = x^T M y$, $M_{ij} = -C_1(a_i, b_j) = C_2(a_i, b_j) \in (-1, 1)$

Thm: Modern proof of minimax

$$1) \max_{x \in S_1} \min_{y \in S_2} x^T M y \leq \min_{y \in S_2} \max_{x \in S_1} x^T M y$$

Oh:

\downarrow For player 1 it is only worse to go first.

second player can make it worse for the first player

but note the first can choose just based on the S_2 choice.

$$2) \max_{x \in S_1} \min_{y \in S_2} x^T M y \geq \min_{y \in S_2} \max_{x \in S_1} x^T M y$$

Oh:

Apply no regret dynamics with average regret ϵ . ($T \geq 4 \log(\max(m, n) / \epsilon^2)$)

For $t=1 \dots T$, let $p^t =$ mixed strategy used by 1
 $q^t =$ 2

Define average strategies $\bar{x} = \frac{1}{T} \sum_{t=1}^T p^t$ for player 1

$\bar{y} = \frac{1}{T} \sum_{t=1}^T q^t$ for player 2

Average cost of 2 = $v = \frac{1}{T} \sum_{t=1}^T (p^t)^T M q^t$

Assume that player 1 plays a_i pure strategy with $p(a_i)=1$

\Rightarrow mixed strategy vector $e_i = (0, \dots, 0, 1, 0, \dots, 0) = -L_i^T$

$$\text{Now payoff of 1 against } \bar{y} = e_i^T M \bar{y} = \frac{1}{T} \sum_{t=1}^T e_i^T M q^t \leq \frac{1}{T} \sum_{t=1}^T (p^+)^T M q^t + \epsilon$$

by definition of \bar{y}

But we can observe: $L_{A_i}^T \leq L_i^T + \epsilon^T$

because average regret $\leq \epsilon$

$\forall x \in S_1$ is convex combination of some pure strategy $e_1 \dots e_m$, also linear, therefore the inequality holds even for mixed strategy vectors.

also: $\Rightarrow \forall x \in S_1: x^T M \bar{y} \leq v + \epsilon$ (*)

By symmetry: $\forall y \in S_2: \bar{x}^T M y \geq v - \epsilon$ (**)

We want to prove $\max_{x \in A_1} \min_{y \in A_2} x^T M y = \min_{y \in S_2} (\bar{x})^T M y \geq$ (***)

$$\stackrel{(**)}{\geq} v - \epsilon \stackrel{(*)}{\geq} \max_{x \in A_1} x^T M \bar{y} - 2\epsilon \geq \min_{y \in S_2} \max_{x \in S_1} x^T M y - 2\epsilon$$

works for any x

arbitr. number chosen.

With $T \rightarrow \infty$, we have $\epsilon \rightarrow 0$ \square

Coarse correlated equilibria (CCE)

Again a prob. distr. p on A , if $\forall i \in A \forall a_i' \in A_i$:

$$\sum_{a \in A} C_i(a) p(a) \leq \sum_{a \in A} C_i(a_i', a_{-i}) p(a)$$

- i does not want to deviate in expectation (before he knows a_{-i} that is suggested to him)

CCE are more general than the true CE.

→ therefore every CE is CCE.

For $\epsilon > 0$, ϵ -CCE is a prob. distr. p on A such that:

$\forall i \in P, \forall a_i' \in A_i$:

$$\sum_{a \in A} C_i(a) p(a) = \sum_{a \in A} C_i(a_i', a_{-i}) p(a) + \epsilon$$

Thm:

$\forall G = (P, A, C), \forall \epsilon > 0, \forall T = T(\epsilon)$:

After T steps of no-regret dynamics, with average regret $\leq \epsilon$,

$$\text{set } p^t = \prod_{i=1}^T p_i^t, \quad p = \frac{1}{T} \sum_{t=1}^T p^t$$

The p is ϵ -CCE.

Oh: We want to show that p is ϵ -CCE.

That is: $\forall i \in P \forall a_i' \in A_i: \mathbb{E}_{a \sim p} [C_i(a)] \leq \mathbb{E}_{a \sim p} [C_i(a_i', a_{-i})] + \epsilon$

definition of p

$$\Rightarrow \mathbb{E}_{a \sim p} [C_i(a)] = \frac{1}{T} \sum_{t=1}^T \mathbb{E}_{a \sim p^t} [C_i(a)]$$

cumulative loss of i in no regret dynamics / T

$$\mathbb{E}_{a \sim p} [C_i(a_i', a_{-i})] = \frac{1}{T} \sum_{t=1}^T \mathbb{E}_{a \sim p^t} [C_i(a_i', a_{-i})]$$

1) → by average regret $\leq \epsilon$

cumulative loss of playing a_i' in no regret dynamics / T

+ ϵ

□